6.1 Lists as an ADT focus on what lists do properties of data homogeneous finite length sequential operations constructor - creates empty list empty - check if list is empty insert - add an item delete - remove an item traverse - various operations that go through a list sequentially search - find an element output - print list contents copy - create a copy sort - rearrange elements various implementations possible static arrays dynamic arrays linked lists w/ pointers (most common) linked lists w/ 2D arrays 6.2 Static Array Implementation list will have max capacity equal to the size of the array list stored sequentially in memory must be able to allocate mem chunk of appropriate size head of list is slot 0 need to add a count of elements in list tells if list is empty or full operations constructor - sets count of elements to 0 empty - check if count is 0 traverse - for (i=0; i<count; i++) takes n for loops to traverse search traversal averages half the list sort traversal can take longer depending on the sort algorithm insert - depends on type of insert tail insert if (count < max-capacity) array[count++] = elementelse // no space left error insert mid-list or head have to move elements down a slot have to validate given position can be used for tail insert too if (count < max-capacity) {

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if (pos < 0 || pos > count)
                 // issue "bad position" error
              else {
                 for(i = count; i > pos; i--)
                    array[i] = array[i-1];
                  array[pos] = element;
                  count++;
              }
            }
           else {
              // issue "no space" error
            }
           takes up to n for loops to shift current elements
              worst case is head insert
              average case is mid-list insert
              best case is tail insert
      delete - also has to shift elements
         if (empty()) {
           // issue "empty list "error
         }
         else if (pos < 0 || pos >= count) {
           // issue "illegal position" error
         }
        else {
           for (i = pos; i < count-1; i++)
              array[i] = array[i + 1];
           count--;
         }
         best/worst/average same as for insert
   implementation details
      how to define element type
         #define macro
        typedef (book)
         template class
      how to define max capacity
         #define macro
         const int (book)
         class variable - should be static
            only one copy of static vars across all class instances
           static const int capacity = 100;
      pp 262-269 has book's implementation
6.3 Dynamic Amery Implementation
   Operations similar to above
      default constructor should select some default capacity &
      allocate mem
      add constructor to take an int for capacity & allocate mem
      add destructor to deallocate mem
      copy has to allocate space for new list first
      add assignment operator to deal w/ memory allocation issue
         otherwise both lists point to same mem
         also could cause mem leaks by not deleting old var
   Implementation changes
      add capacity var to member vars
```

add capacity var to member vars change array var to a pointer 6.4 Linked Lists use pointers to connect elements arrays have implicit order linked lists have explicit order list nodes need to store data & point to the next element create node as separate class needs functions to retrieve/set data & retrieve/set pointer list is a collection of nodes & operations on the nodes needs a pointer the 1st (head) node consider list w/ only head ptr now list variants add other pointers list operations create empty list - set head to NULL is empty? - does head equal NULL? traversal - from head node, follow pointer to next element repeat until pointer to next is NULL pseudocode set ptr to head while ptr is not NULL do traversal operation set ptr to ptr's next node insertion - add new node to list several cases depending on where adding head insert / 1st node insert new node will become head pseudocode set new's next to head set head to new must set next before changing head pointer otherwise lose reference to old list tail insert new node will become end of list pseudocode traverse list to find current tail set tail's next to new set new's next to NULL mid-list insert insert after some specific node pseudocode traverse list to find previous node set new's next to prev's next set prev's next to new can avoid traversal if pass ptr to insert pseudocode if ptr is NULL, do head insert else, do tail/mid-list insert w/ ptr as tail/previous traversal still has to be done somewhere for mid-list/tail insert to make insertion of 1st element or at tail easier, have list node initialize next to NULL deletion - remove node from list must update pointers to reflect new order

two cases head delete remove 1st element 2nd element becomes new head Pseudocode create tmp ptr that points to head set head to head's next deallocate tmp mid-list & tail delete need to find nod before one being deleted previous node will "skip over" deleted node pseudocode traverse list to find previous node set prev's next to node's next deallocate node 6.5 Linked List Implementation NOTE: This differs from the book's implementation Node class member vars an element (template type) a pointer to the next node member functions default constructor - sets next to NULL a constructor that takes an element & sets next to NULL a constructor that takes an element & node pointer setData to set the element getData to retrieve element setNext to set next pointer getNext to retrieve next pointer equality operator (for list search) output operator (for list output) Linked List Class member vars a node pointer for head member functions default constructor - sets head to NULL destructor - deallocate list nodes copy constructor - create 2nd list that stores same elements has separate pointers & memory space assignment operator - also create copy bool empty() - check if head is NULL output operator - print list contents Node \*search (T elem) - search list for element return pointer of node if found return NULL if not found Node \*find previous (Node \*ptr) find the node before given node can be private helper function for delete void insert (T elem, Node \*prev) create new node to store elem insert at head if prev is NULL otherwise insert after prev void delete (Node \*node) - remove node from list

void delete (T elem) - alt form of delete traverse list to find elem's node call node delete function